## Exercise 2.1: /10

Let A be a non-c.e. set. Use the finite extension method to prove that there exists a set B which is not A-c.e. and such that A is not B-c.e.

## Solution 2.1:

Recall that  $W_e^X = \operatorname{dom} \Phi_e^X$ . We want to satisfy the following requirements for every  $e \in \mathbb{N}$ :

- $\mathcal{R}_e$ :  $W_e^A \neq B$ .
- $\mathcal{S}_e: W_e^B \neq A.$

If all the  $\mathcal{R}$ -requirements are satisfied, then B is not A-c.e. If all the  $\mathcal{S}$ -requirements are satisfied, then A is not B-c.e.

Satisfying  $\mathcal{R}_e$ . Assume  $\sigma_n$  is defined. Let  $x = |\sigma_n|$ . If  $\Phi_e^A(x) \uparrow$ , then take  $\sigma_{n+1} = \sigma_n 1$ . Otherwise, let  $\sigma_{n+1} = \sigma_n 0$ .

Satisfying  $S_e$ . Assume  $\sigma_n$  is defined. We have three cases:

- Case 1: there is some  $k \notin A$  and some  $\tau \succeq \sigma_n$  such that  $k \in W_e^{\tau}$ . Then let  $\sigma_{n+1} = \tau$ .
- Case 2: there is some  $k \in A$  such that for every  $\tau \succeq \sigma_n, k \notin W_e^{\tau}$ . Then let  $\sigma_{n+1} = \sigma_n$ .
- Case 3: none of the cases above hold. Then  $A = \{k : \exists \tau \succeq \sigma_n k \in W_e^{\tau}\}$  so A is c.e, contradiction.

Recall that  $W_e = \operatorname{dom} \Phi_e$ .

## Exercise 2.2: /10

Show that  $Tot = \{e : W_e = \mathbb{N}\}$  is  $\Pi_2^0$ -complete.

## Solution 2.2:

First, let us show that Tot is  $\Pi_2^0$ .

$$\mathsf{Tot} = \{ e : \forall x \exists t \Phi_e(x)[t] \downarrow \}$$

The predicate  $P(e, x, t) \equiv \Phi_e(x)[t] \downarrow$  being computable, Tot is  $\Pi_2^0$ .

Let  $A = \{x : \forall y \exists z R(x, y, z)\}$  be a  $\Pi_2^0$  set. Let us show that  $A \leq_m$  Tot. Let  $f : \mathbb{N} \to \mathbb{N}$ be the computable function which to every n, associates the code f(n) of the function  $\Phi_{f(n)}$ defined as follows: On input y,  $\Phi_{f(n)}$  searches for some z such that R(n, y, z) holds. If found,  $\Phi_{f(n)}(y) \downarrow$ , otherwise  $\Phi_{f(n)}(y) \uparrow$ .

- If n ∈ A, then for every y, there is some z such that R(n, y, z) holds, so Φ<sub>f(n)</sub> is total, so f(n) ∈ Tot.
- If  $n \notin A$ , then there is some y such that for every z, R(n, y, z) does not hold, so  $\Phi_{f(n)}(y) \uparrow$ , hence  $f(n) \notin \text{Tot.}$