

Iterative Forcing and Hyperimmunity in Reverse Mathematics

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Abstract. The separation between two theorems in reverse mathematics is usually done by constructing a Turing ideal satisfying a theorem P and avoiding the solutions to a fixed instance of a theorem Q. Lerman, Solomon and Towsner introduced a forcing technique for iterating a computable non-reducibility in order to separate theorems over omega-models. In this paper, we present a modularized version of their framework in terms of preservation of hyperimmunity and show that it is powerful enough to obtain the same separations results as Wang did with his notion of preservation of definitions.

1 Introduction

Reverse mathematics is a mathematical program which aims to capture the provability content of ordinary (i.e. non set-theoretic) theorems. It uses the framework of subsystems of second-order arithmetic, with a base theory RCA_0 which is composed of the basic axioms of Peano arithmetic together with the Δ_1^0 comprehension scheme and the Σ_1^0 induction scheme. Thanks to the equivalence between Δ_1^0 -definable sets and computable sets, RCA_0 can be thought as capturing “computational mathematics”. See [8] for a good introduction.

Many theorems are Π_2^1 statements $(\forall X)(\exists Y)\Phi(X, Y)$ and come with a natural class of *instances* X . The sets Y such that $\Phi(X, Y)$ holds are *solutions* to X . For example, König’s lemma (KL) states that every infinite, finitely branching tree has an infinite path. An instance of KL is an infinite, finitely branching tree T . A solution to T is an infinite path through T . Given two Π_2^1 statements P and Q, proving an implication $Q \rightarrow P$ over RCA_0 consists of taking a P-instance X and constructing a solution to X through a computational process involving several applications of the Q statement. Empirically, many proofs of implications are in fact *computable reductions* [9].

Definition 1 (Computable reducibility). Fix two Π_2^1 statements P and Q. We say that P is *computably reducible* to Q (written $P \leq_c Q$) if every P-instance I computes a Q-instance J such that for every solution X to J , $X \oplus I$ computes a solution to I .

If the computable reduction between from P to Q can be formalized over RCA_0 , then $\text{RCA}_0 \vdash Q \rightarrow P$. However, P may not be computably reducible to Q while $\text{RCA}_0 \vdash Q \rightarrow P$. Indeed, one may need more than one application of Q to solve the instance

of P. This is for example the case of Ramsey's theorem for pairs with n colors (RT_n^2) which implies RT_{n+1}^2 over RCA_0 , but $RT_{n+1}^2 \not\leq_c RT_n^2$ for $n \geq 1$ (see [21]).

In order to prove the non-implication between P and Q, one needs to iterate the computable non-reducibility in order to build a model of Q which is not a model of P. This is the purpose of the framework developed by Lerman, Solomon and Towsner in [14]. They successfully used their framework for separating the Erdős-Moser theorem (EM) from the stable ascending descending sequence principle (SADS) and separating the ascending descending sequence (ADS) from the stable chain antichain principle (SCAC). Their approach has been reused by Flood & Towsner [5] and the author [19] on diagonal non-computability statements.

However, their framework suffers some drawbacks. In particular the forcing notions involved are heavy and the deep combinatorics witnessing the non-implications are hidden by the complexity of the proof. Moreover, the P-instance chosen in the ground forcing depends on the forcing notion used in the iteration forcing and therefore the overall construction is not modular. On the other hand, Wang [23] recently introduced the notion of preservation of definitions and made independent proofs of preservations for various statements included EM. Then he deduced that the conjunction of those statements does not imply SADS, therefore strengthening the result of Lerman, Solomon & Towsner in a modular way. Variants of this notion have been reused by the author [21] for separating the free set theorem (FS) from RT_2^2 .

In this paper, we present a modularized version of the framework of Lerman, Solomon & Towsner and use it to reprove the separation results obtained by Wang [23]. We thereby show that this framework is a viable alternative to the notion introduced by Wang for separating statements in reverse mathematics. In particular, we reprove the following theorem, in which COH is the cohesiveness principle, WKL_0 is weak König's lemma, RRT_2^2 the rainbow Ramsey theorem for pairs, $\Pi_1^0 G$ the Π_1^0 -genericity principle and STS^2 the stable thin set theorem for pairs.

Theorem 2 (Wang [23]). *Let Φ be the conjunction of COH, WKL_0 , RRT_2^2 , $\Pi_1^0 G$, and EM. Over RCA_0 , Φ does not imply any of SADS and STS^2 .*

In section 2, we introduce the framework of Lerman, Solomon & Towsner in its original form and detail its drawbacks. Then, in section 3, we develop a modularized version of their framework. In section 4, we establish basic preservation results, before reproving in section 5 Wang's theorem. Last, we reprove in section 6 the separation obtained by the author in [21].

1.1 Notation

String, sequence. Fix an integer $k \in \omega$. A *string* (over k) is an ordered tuple of integers a_0, \dots, a_{n-1} (such that $a_i < k$ for every $i < n$). The empty string is written ε . A *sequence* (over k) is an infinite listing of integers a_0, a_1, \dots (such that $a_i < k$ for every $i \in \omega$). Given $s \in \omega$, k^s is the set of strings of length s over k and $k^{<\omega}$ is the set of finite strings over k . Given a string $\sigma \in k^{<\omega}$, we denote by $|\sigma|$ its length. Given two strings $\sigma, \tau \in k^{<\omega}$, σ is a *prefix* of τ (written $\sigma \preceq \tau$) if there exists a string $\rho \in k^{<\omega}$ such that $\sigma\rho = \tau$. A *binary string* (resp. real) is a *string* (resp. sequence)

over 2. We may equate a real with a set of integers by considering that the real is its characteristic function.

Tree, path. A tree $T \subseteq \omega^{<\omega}$ is a set downward-closed under the prefix relation. The tree T is *finitely branching* if every node $\sigma \in T$ has finitely many immediate successors. A *binary tree* is a tree $T \subseteq 2^{<\omega}$. A set $P \subseteq \omega$ is a *path* through T if for every $\sigma \prec P$, $\sigma \in T$. A string $\sigma \in k^{<\omega}$ is a *stem* of a tree T if every $\tau \in T$ is comparable with σ . Given a tree T and a string $\sigma \in T$, we denote by $T^{[\sigma]}$ the subtree $\{\tau \in T : \tau \preceq \sigma \vee \tau \succeq \sigma\}$.

Sets. Given two sets X and Y , $X \subseteq^* Y$ means that X is almost included into Y , $X =^* Y$ means $X \subseteq^* Y \wedge Y \subseteq^* X$ and $X \subseteq_{\text{fin}} Y$ means that X is a finite subset of Y . Given some $x \in \omega$, $A > x$ denotes the formula $(\forall y \in A)[y > x]$.

2 The Iteration Framework

An ω -structure is a structure $\mathcal{M} = (\omega, S, +, \cdot, <)$ where ω is the set of standard integers, $+$, \cdot and $<$ are the standard operations over integers and S is a set of reals such that \mathcal{M} satisfies the axioms of RCA_0 . Friedman [7] characterized the second-order parts S of ω -structures as those forming a *Turing ideal*, that is, a set of reals closed under Turing join and downward-closed under the Turing reduction.

Fix two Π_2^1 statements P and Q . The construction of an ω -model of P which is not a model of Q consists of creating a Turing ideal \mathcal{I} together with a fixed Q -instance $I \in \mathcal{I}$, such that every P -instance $J \in \mathcal{I}$ has a solution in \mathcal{I} , whereas I contains no solution in \mathcal{I} . In the first place, let us just focus on the one-step case, that is, a proof that $Q \not\leq_c P$. To do so, one has to choose carefully some Q -instance I such that every I -computable P -instance has a solution X which does not I -compute a solution to I . The construction of a solution X to some I -computable P -instance J will have to satisfy the following scheme of requirements for each index e :

$$\mathcal{R}_e : \Phi_e^{X \oplus I} \text{ infinite} \rightarrow \Phi_e^{X \oplus I} \text{ is not a solution to } I$$

Such requirements may not be satisfiable for an arbitrary Q -instance I . The choice of the instance and the satisfaction of the requirement is strongly dependent on the combinatorics of the statement Q and the forcing notion used for constructing a solution to J . A recurrent approach in the framework of Lerman, Solomon & Towsner consists of constructing a Q -instance I which satisfies some fairness property. The forcing notion \mathbb{P}^I used in the construction of a solution to J is usually designed so that

- (i) There exists an I -computable set encoding (at least) every condition in \mathbb{P}^I
- (ii) Given some forcing condition in \mathbb{P}^I , one can uniformly find in a c.e. search a finite set of candidate extensions such that one of them is in \mathbb{P}^I (e.g. the notion of split pair in [14], of finite cover for a tree forcing, ...).

The fairness property states the following:

“For every condition in \mathbb{P}^I , if for every $x \in \omega$, there exists a *finite* Q -instance $A > x$ and a finite set of candidate extensions d_0, \dots, d_m such that $\Phi_e^{d_i \oplus I}$ is not a solution to A for each $i \leq m$, then one of the A 's is a subinstance of I .”

This property is designed so that we can satisfy it by taking each condition $c \in \mathbb{P}^I$ one at a time, find some finite Q-instance A on which I is not yet defined, and define I over A . One can think of the instance I as a fair adversary who, if we have infinitely often the occasion to beat him, will be actually beaten at some time.

Suppose now we want to extend this computable non-reducibility into a separation over ω -structures. One may naturally try to make the instance I satisfy the fairness property at every level of the iteration forcing. At the first iteration with an I -computable P-instance J , the property is unchanged. At the second iteration, the P-instance J_1 is $X_0 \oplus I$ -computable, but the set X_0 is not yet constructed. Thankfully, the fairness property requires a finite piece of oracle X_0 . Therefore we can modify the fairness property which becomes

“For every condition $c_0 \in \mathbb{P}^I$ and every condition $c_1 \in \mathbb{P}^{c_0 \oplus I}$, if for every $x \in \omega$, there exists a Q-instance $A > x$, a finite set of candidate extensions $d_0, \dots, d_m \in \mathbb{P}^I$ and $d_{0,i}, \dots, d_{n_i,i} \in \mathbb{P}^{d_i \oplus I}$ for each $i \leq m$ such that $\Phi_e^{d_{j,i} \oplus d_i \oplus I}$ is not a solution to A for each $i \leq m$ and $j \leq n_i$, then one of the A 's is a subinstance of I .”

Since this property becomes overly complicated in the general case, Lerman, Solomon and Towsner abstracted the notion of requirement and made it a $\Sigma_1^{0,I}$ black box which takes as parameters a condition and a finite Q-instance. Instead of making the instance I in charge of satisfying the fairness property at every level of the iteration forcing, the instance I satisfies the property only at the first level. Then, by encoding a requirement at the next level into a requirement at the current level, the iteration forcing ensures the propagation of this fairness property from the first level to every level. The property in its abstracted form is then

“For every condition in \mathbb{P}^I and every $\Sigma_1^{0,I}$ predicate \mathcal{K}^I , if for every $x \in \omega$, there is a *finite* Q-instance $A > x$ and a finite set of candidate extensions d_0, \dots, d_m such that $\mathcal{K}^I(A, d_i)$ is satisfied for each $i \leq m$, then one of the A 's is a subinstance of I .”

In particular, by letting $\mathcal{K}^I(A, c)$ be the predicate “ $\Phi_e^{d_i \oplus I}$ is not a solution to A ”, the requirements \mathcal{R}_e will be satisfied.

The problem of such an approach is that the construction of the Q-instance strongly depends on the forcing notion used in the iteration forcing. A slight modification of the latter requires to modify the ground forcing. Moreover, if someone wants to prove that the conjunction of two statements does not imply a third one, we need to construct an instance I which will satisfy the fairness property for the two statements, and in each iteration forcing, we will need to ensure that both properties are propagated to the next iteration. The size of the overall construction explodes when trying to make a separation of the conjunction of several statements at the same time.

3 Preservation of Hyperimmunity

In this section, we propose a general simplification of the framework of Lerman, Solomon & Towsner [14] and illustrate it in the case of the separation of EM from SADS. The corresponding fairness property happens to coincide with the notion of hyperimmunity. The underlying idea ruling this simplification is the following: since each

condition in the iteration forcing can be given an index and since the finite set of candidate extensions of a condition c , can be found in a c.e. search, given a $\Sigma_1^{0,I}$ predicate \mathcal{K}^I , the following formula is again $\Sigma_1^{0,I}$:

$\varphi(U) = \text{“there exists a finite set of candidate extensions } d_0, \dots, d_m \text{ of } c \text{ such that } \mathcal{K}^I(U, d_i) \text{ is satisfied for each } i \leq m\text{”}$

We can therefore abstract the iteration forcing and ask the instance I to satisfy the following property:

“For every $\Sigma_1^{0,I}$ predicate $\varphi(U)$, if for every $x \in \omega$, there exists a finite Q-instance $A > x$ such that $\varphi(A)$ is satisfied, then one of the A 's is a subinstance of I .”

Let us illustrate how this simplification works by reproving the separation of the Erdős-Moser theorem from the ascending descending sequence principle.

Definition 3 (Ascending descending sequence). *A linear order is stable if it is of order type $\omega + \omega^*$. ADS is the statement “Every linear order admits an infinite ascending or descending sequence”. SADS is the restriction of ADS to stable linear orders.*

The ascending descending sequence principle has been studied within the framework of reverse mathematics by Hirschfeldt & Shore [10]. Lerman, Solomon & Towsner [14] constructed an infinite stable linear order I with ω and ω^* parts respectively B_0 and B_1 , such that for every condition c and every $\Sigma_1^{0,I}$ predicate \mathcal{K}^I , if for every $x \in \omega$, there exists a finite set $A > x$ and a finite set of candidate extensions d_0, \dots, d_m of c such that $\mathcal{K}^I(A, d_i)$ is satisfied for each $i \leq m$, then one of the A 's will be included in B_0 and another one will be included in B_1 . In particular, taking $\mathcal{K}^I(A, c) = \Phi_e^{c \oplus I} \cap A \neq \emptyset$, no infinite solution to the constructed tournament I -computes a solution to I . After abstraction, we obtain the following property:

“For every $\Sigma_1^{0,I}$ predicate $\varphi(U)$, if for every $x \in \omega$, there exists a finite set $A > x$ such that $\varphi(A)$ is satisfied, one of the A 's is included in B_0 and one of the A 's is included in B_1 .”

Following the terminology of [14], we say that a formula $\varphi(U)$ is *essential* if for every $x \in \omega$, there exists some finite set $A > x$ such that $\varphi(A)$ holds. This fairness property coincides with the notion of hyperimmunity for $\overline{B_0}$ and $\overline{B_1}$.

Definition 4 (Preservation of hyperimmunity).

1. Let D_0, D_1, \dots be a computable list of all finite sets and let f be computable. A c.e. array $\{D_{f(i)}\}_{i \geq 0}$ is a c.e. set of mutually disjoint finite sets $D_{f(i)}$. A set B is hyperimmune if for every c.e. array $\{D_{f(i)}\}_{i \geq 0}$, $D_{f(i)} \cap B = \emptyset$ for some i .
2. A Π_2^1 statement P admits preservation of hyperimmunity if for each set Z , each countable collection of Z -hyperimmune sets A_0, A_1, \dots , and each P -instance $X \leq_T Z$ there exists a solution Y to X such that the A 's are $Y \oplus Z$ -hyperimmune.

The following lemma establishes the link between the fairness property for SADS and the notion of hyperimmunity.

Lemma 5. *Fix a set Z . A set B is Z -hyperimmune if and only if for every essential $\Sigma_1^{0,Z}$ predicate $\varphi(U)$, $\varphi(A)$ holds for some finite set $A \subseteq \overline{B}$.*

Hirschfeldt, Shore & Slaman constructed in [11, Theorem 4.1] a stable computable linear order such that both the ω and the ω^* part are hyperimmune. As every ascending (resp. descending) sequence is an infinite subset of the ω (resp. ω^*) part of the linear order, we deduce the following theorem.

Theorem 6. *SADS does not admit preservation of hyperimmunity.*

A slight modification of the forcing in [14] gives preservation of hyperimmunity of the Erdős-Moser theorem. We will however reprove it in a later section with a simpler forcing notion. As expected, the notion of preservation of hyperimmunity can be used to separate statements in reverse mathematics.

Lemma 7. *Fix two Π_2^1 statements P and Q. If P admits preservation of hyperimmunity and Q does not, then P does not imply Q over RCA_0 .*

Before starting an analysis of preservations of hyperimmunity for basic statements, we state another negative preservation result which enables to reprove that the Erdős-Moser theorem does not imply the stable thin set theorem for pairs [15].

Definition 8 (Thin set theorem). *Let $k \in \omega$ and $f : [\omega]^k \rightarrow \omega$. A set A is f-thin if $f([A]^n) \neq \omega$, that is, if the set A “avoids” at least one color. TS^k is the statement “every function $f : [\omega]^k \rightarrow \omega$ has an infinite f-thin set”. STS^2 is the restriction of TS^2 to stable functions.*

Introduced by Friedman in [6], the basic reverse mathematics of the thin set theorem has been settled by Cholak, Hirst & Jockusch in [2]. Its study has been continued by Wang [24], Rice [22] and the author [17,21]. The author constructed in [20] an infinite computable stable function $f : [\omega]^2 \rightarrow \omega$ such that the sets $B_i = \{n \in \omega : \lim_s f(n, s) \neq i\}$ are all hyperimmune. Every infinite f-thin set being an infinite subset of one of the B's, we deduce the following theorem.

Theorem 9. *STS^2 does not admit preservation of hyperimmunity.*

4 Basic Preservations of Hyperimmunity

When defining a notion, it is usually convenient to see how it relates with typical sets. There are two kinds of typicalities: genericity and randomness. Both notions admit preservation of hyperimmunity.

Theorem 10. *Fix some set Z and a countable collection of Z-hyperimmune sets B_0, B_1, \dots*

1. *If G is sufficiently Cohen generic relative to Z, the B's are $G \oplus Z$ -hyperimmune.*
2. *If R is sufficiently random relative to Z, the B's are $R \oplus Z$ -hyperimmune.*

Note that this does not mean that the sets G and R are hyperimmune-free relative to Z. In fact, the converse holds: if G is sufficiently generic and R sufficiently random, then both are Z-hyperimmune. Some statements like the atomic model theorem (AMT), Π_1^0 -genericity (Π_1^0G) and the rainbow Ramsey theorem for pairs (RRT_2^2)

are direct consequences of genericity and randomness [11,4]. We can deduce from Theorem 10 that they all admit preservation of hyperimmunity.

Cohesiveness is a very useful statement in the analysis of Ramsey-type theorems as it enables one to transform an arbitrary instance into a stable one [3]. A set C is *cohesive* for a sequence of sets R_0, R_1, \dots if $C \subseteq^* R_i$ or $C \subseteq^* \overline{R_i}$ for each i .

Theorem 11. *COH admits preservation of hyperimmunity.*

The proof is done by the usual construction of a cohesive set with Mathias forcing, combined with the following lemma.

Lemma 12. *For every set Z , every Z -computable Mathias condition (F, X) , every $\Sigma_1^{0,Z}$ formula $\varphi(G, U)$ and every Z -hyperimmune set B , there exists an extension (E, Y) such that $X =^* Y$ and either $\varphi(G, U)$ is not essential for every set G satisfying (E, Y) , or $\varphi(E, A)$ holds for some finite set $A \subseteq \overline{B}$.*

Proof. Define

$$\psi(U) = (\exists G \supseteq F)[G \subseteq F \cup X \wedge \varphi(G, U)]$$

The formula $\psi(U)$ is $\Sigma_1^{0,Z}$. By hyperimmunity of B , either $\psi(U)$ is not essential, or $\psi(A)$ holds for some finite set $A \subseteq \overline{B}$. In the first case, the condition (F, X) already satisfies the desired property. In the second case, let $A \subseteq_{\text{fin}} \overline{B}$ be such that $\psi(A)$ holds. By the use property, there exists a finite set E satisfying (F, X) such that $\varphi(E, A)$ holds. Let $Y = X \setminus [0, \max(E)]$. The condition (E, Y) is a valid extension. \square

Weak König's lemma (WKL_0) states that every infinite, binary tree admits an infinite path.

Theorem 13. *WKL_0 admits preservation of hyperimmunity.*

Wei Wang [personal communication] observed that WKL_0 preserves hyperimmunity in a much stronger sense than COH, since cohesive sets are of hyperimmune degree [12], whereas by the hyperimmune-free basis theorem [13], WKL_0 can preserve hyperimmunities of every hyperimmune set simultaneously and not only countably many.

5 The Erdős-Moser Theorem and Preservation of Hyperimmunity

The Erdős-Moser theorem is a statement from graph theory which received a particular interest from reverse mathematical community as it provides, together with the ascending descending sequence principle, an alternative proof of Ramsey's theorem for pairs.

Definition 14 (Erdős-Moser theorem). *A tournament T is an irreflexive binary relation such that for all $x, y \in \omega$ with $x \neq y$, exactly one of $T(x, y)$ or $T(y, x)$ holds. A tournament T is transitive if the corresponding relation T is transitive in the usual sense. EM is the statement "Every infinite tournament T has an infinite transitive subtournament."*

The Erdős-Moser theorem was introduced in reverse mathematics by Bovykin & Weiermann [1] and then studied by Lerman, Solomon & Towsner [14] and the author [18,17,16]. In this section, we give a simple proof of the following theorem.

Theorem 15. *EM admits preservation of hyperimmunity.*

The proof of Theorem 15 exploits the modularity of the framework by using preservation of hyperimmunity of WKL_0 . Together with the previous preservation results, this theorem is sufficient to reprove Theorem 2. We must first introduce some terminology.

Definition 16 (Minimal interval). *Let T be an infinite tournament and $a, b \in T$ be such that $T(a, b)$ holds. The interval (a, b) is the set of all $x \in T$ such that $T(a, x)$ and $T(x, b)$ hold. Let $F \subseteq T$ be a finite transitive subtournament of T . For $a, b \in F$ such that $T(a, b)$ holds, we say that (a, b) is a minimal interval of F if there is no $c \in F \cap (a, b)$, i.e., no $c \in F$ such that $T(a, c)$ and $T(c, b)$ both hold.*

Definition 17. *An Erdős Moser condition (EM condition) for an infinite tournament T is a Mathias condition (F, X) where*

- (a) $F \cup \{x\}$ is T -transitive for each $x \in X$
- (b) X is included in a minimal T -interval of F .

EM extension is Mathias extension. A set G satisfies an EM condition (F, X) if it is T -transitive and satisfies the Mathias extension (F, X) . Basic properties of EM conditions have been stated and proven in [18].

Fix a set Z and some countable collection of Z -hyperimmune sets B_0, B_1, \dots . Our forcing notion is the partial order of Erdős Moser conditions (F, X) such that the B 's are $X \oplus Z$ -hyperimmune. Our initial condition is (\emptyset, ω) . By Lemma 5.9 in [18], EM conditions are extendable, so we can force the transitive subtournament to be infinite. Therefore it suffices to prove the following lemma to deduce Theorem 15.

Lemma 18. *Fix a condition (F, X) , some $i \in \omega$ and some $\Sigma_1^{0,Z}$ formula $\varphi(G, U)$. There exists an extension (E, Y) such that either $\varphi(G, U)$ is not essential for every set G satisfying (E, Y) , or $\varphi(E, A)$ holds for some finite set $A \subseteq \overline{B_i}$.*

Proof. Let $\psi(U)$ be the formula “For every partition $X_0 \cup X_1 = X$, there exists some $j < 2$, a T -transitive set $G \subseteq X_j$ and a set $\tilde{A} \subseteq U$ such that $\varphi(F \cup G, \tilde{A})$ holds.” By compactness, $\psi(U)$ is a $\Sigma_1^{0, X \oplus Z}$ formula. By $X \oplus Z$ -hyperimmunity of B_i , we have two cases:

- Case 1: $\psi(A)$ holds for some finite set $A \subseteq \overline{B_i}$. By compactness, there exists a finite set $H \subset X$ such that for every partition $H_0 \cup H_1 = H$, there exists some $j < 2$, a T -transitive set $G \subseteq H_j$ and a set $\tilde{A} \subseteq A$ such that $\varphi(F \cup G, \tilde{A})$ holds. Given two sets U and V , we denote by $U \rightarrow_T V$ the formula $(\forall x \in U)(\forall y \in V)T(x, y)$. Each element $y \in X$ induces a partition $H_0 \cup H_1 = H$ such that $H_0 \rightarrow_T \{y\} \rightarrow_T H_1$. There exists finitely many such partitions, so by the infinite pigeonhole principle, there exists an X -computable infinite set $Y \subset X$ and a partition $H_0 \cup H_1 = H$ such that $H_0 \rightarrow_T Y \rightarrow_T H_1$. Let $j < 2$ and $G \subseteq H_j$ be the T -transitive set such that $\varphi(F \cup G, \tilde{A})$ holds for some $\tilde{A} \subseteq A \subseteq \overline{B_i}$. By Lemma 5.9 in [18], $(F \cup G, Y)$ is a valid extension.

- Case 2: $\psi(U)$ is not essential with some witness x . Then the $\Pi_1^{0,X \oplus Z}$ class \mathcal{C} of sets $X_0 \oplus X_1$ such that $X_0 \cup X_1 = X$ and for every $j < 2$, every T -transitive set $G \subseteq X_j$ and every finite set $\tilde{A} \succ x$, the formula $\varphi(F \cup G, \tilde{A})$ does not hold is not empty. By preservation of hyperimmunity of WKL_0 , there exists some partition $X_0 \oplus X_1 \in \mathcal{C}$ such that the B 's are $X_0 \oplus X_1 \oplus Z$ -hyperimmune. The set X_j is infinite for some $j < 2$ and the condition (F, X_i) is the desired EM extension. \square

6 Thin Set Theorem and Preservation of Hyperimmunity

There exists a fundamental difference in the way SADS and STS^2 witness their failure of preservation of hyperimmunity. In the case of SADS, we construct two hyperimmune sets whereas in the case of STS^2 , a countable collection of hyperimmune sets is used. This difference can be exploited to obtain further separation results.

Definition 19 (Preservation of n hyperimmunities). A Π_2^1 statement P admits preservation of n hyperimmunities if for each set Z , each Z -hyperimmune sets A_0, \dots, A_{n-1} , and each P -instance $X \leq_T Z$ there exists a solution Y to X such that the A 's are $Y \oplus Z$ -hyperimmune.

Theorem 6 shows that SADS does not admit preservation of 2 hyperimmunities. On the other hand, we shall see that STS^2 admits preservation of n hyperimmunities for every $n \in \omega$. Consider the following variants of the thin set theorem.

Definition 20 (Thin set theorem). Given a function $f : [\omega]^k \rightarrow n$, an infinite set H is f -thin if $|f([H]^k)| \leq n-1$ (i.e. f avoids one color over H). For every $k \geq 1$ and $n \geq 2$, TS_n^k is the statement “Every function $f : [\omega]^k \rightarrow n$ has an infinite f -thin set”. STS_n^2 is the restriction of TS_n^2 to stable colorings.

Note that TS_2^2 is Ramsey’s theorem for pairs. The following theorem is sufficient to separate TS^2 from Ramsey’s theorem for pairs as $\text{TS}^2 \leq_c \text{TS}_n^2$ for every $n \geq 2$. The proof of preservation is rather technical and is therefore proven in appendix.

Theorem 21. For every $n \geq 1$, TS_{n+1}^2 admits preservation of n but not $n+1$ hyperimmunities.

In the case $n = 1$, noticing that the arithmetical comprehension scheme (ACA_0) does not preserve 1 hyperimmunities as witnessed by taking any Δ_2^0 hyperimmune set, we re-obtain the separation of Ramsey’s theorem for pairs from ACA_0 . Hirschfeldt & Jockusch [9] asked whether TS_{n+1}^2 implies TS_n^2 over RCA_0 . The author answered negatively in [21]. Preservation of n hyperimmunities gives the same separation.

Theorem 22 (Patey [21]). For every $n \geq 2$, let Φ be the conjunction of COH , WKL_0 , RRT_2^2 , $\Pi_1^0\text{G}$, EM , TS_{n+1}^2 . Over RCA_0 , Φ does not imply any of SADS and STS_n^2 .

Acknowledgements. The author is thankful to Wei Wang for useful comments and discussions.

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A Preservation of hyperimmunity

Proof (Lemma 5). Let D_0, D_1, \dots be a computable list of all finite sets.

- Fix some set Z and some Z -hyperimmune set B . For every essential $\Sigma_1^{0,Z}$ formula $\varphi(U)$, define the Z -computable function f inductively so that $\varphi(D_{f(0)})$ holds and for every i , $D_{f(i+1)} > D_{f(i)}$ and $\varphi(D_{f(i+1)})$ holds. Because $\varphi(U)$ is essential, the function f is total. $\{D_{f(i)}\}_{i \geq 0}$ is a Z -c.e. weak array, so by Z -hyperimmunity, $D_{f(i)} \cap B = \emptyset$ for some i , hence $D_{f(i)} \subseteq \overline{B}$ and $\varphi(D_{f(i)})$ holds.
- Fix some set Z and some set B such that the fairness property of Lemma 5 holds. For every Z -c.e. weak array $\{D_{f(i)}\}_{i \geq 0}$, define the $\Sigma_1^{0,Z}$ formula $\varphi(U) = (\exists i)[U = D_{f(i)}]$. The formula $\varphi(U)$ is essential, so there exists some finite set $A \subseteq \overline{B}$ such that $\varphi(A)$ holds. In particular, there exists some i such that $D_{f(i)} \subseteq \overline{B}$. \square

Proof (Lemma 7). Fix a set X_0 , a countable collection of X_0 -hyperimmune sets B_0, B_1, \dots and an X_0 -computable Q -instance J such that for every solution Y to J , one of the B 's is not $Y \oplus X_0$ -hyperimmune. By preservation of hyperimmunity of P and carefully choosing a sequence of P -instances I_0, I_1, \dots , we can define an infinite sequence of sets X_1, X_2, \dots such that for each $n \in \omega$

- (a) X_{n+1} is a solution to the P -instance $I_n^{X_0 \oplus \dots \oplus X_n}$
- (b) The B 's are $X_0 \oplus \dots \oplus X_n$ -hyperimmune
- (c) For every $X_0 \oplus \dots \oplus X_n$ -computable P -instance I , there exists some m such that $I = I_m^{X_0 \oplus \dots \oplus X_m}$.

Let \mathcal{M} be the ω -structure whose second-order part is the Turing ideal

$$\mathcal{S} = \{Y : (\exists n)[Y \leq_T X_0 \oplus \dots \oplus X_n]\}$$

In particular, the Q -instance J is in \mathcal{S} , but the B 's are Y -hyperimmune for every $Y \in \mathcal{S}$, so J has no solution $Y \in \mathcal{S}$ and $\mathcal{M} \not\models Q$. By construction of \mathcal{S} , every P -instance $I \in \mathcal{S}$ has a solution $X_n \in \mathcal{S}$, so by Friedman [7], $\mathcal{M} \models \text{RCA}_0 \wedge P$. \square

B Basic preservation proofs

The proofs in this appendix are inspired by Wang's work in [23].

Proof (Item 1 of Theorem 10). It suffices to prove that for every $\Sigma_1^{0,Z}$ formula $\varphi(G, U)$ and every $i \in \omega$, the set of conditions σ forcing $\varphi(G, U)$ not to be essential or such that $\varphi(\sigma, A)$ holds for some finite set $A \subset \overline{B_i}$ is dense. Fix any string $\sigma \in 2^{<\omega}$. Define

$$\psi(U) = (\exists \tau \succeq \sigma) \varphi(\tau, U)$$

The formula $\psi(U)$ is $\Sigma_1^{0,Z}$, so by Z -hyperimmunity of B_i , either $\psi(U)$ is not essential, or $\psi(A)$ holds for some finite set $A \subset \overline{B_i}$. If $\psi(U)$ is not essential with witness $x \in \omega$, then σ forces $\varphi(G, U)$ not to be essential with the same witness. If $\psi(U)$ is essential, then there exists some finite set $A \subset \overline{B_i}$ such that $\psi(A)$ holds. Unfolding the definition of $\psi(A)$, there exists some $\tau \succeq \sigma$ such that $\varphi(\tau, A)$ holds. The condition τ is an extension such that $\varphi(\tau, A)$ holds for some $A \subset \overline{B_i}$. \square

Proof (Item 2 of Theorem 10). It suffices to prove that for every $\Sigma_1^{0,Z}$ formula $\varphi(G, U)$ and every $i \in \omega$, the following class is Lebesgue null.

$$\mathcal{S} = \{X : [\varphi(X, U) \text{ is essential}] \wedge (\forall A \subseteq_{\text{fin}} \omega) \varphi(X, A) \rightarrow A \not\subseteq \overline{B_i}\}$$

Suppose it is not the case. There exists $\sigma \in 2^{<\omega}$ such that

$$\mu(X \in \mathcal{S} : \sigma \prec X) > 2^{-|\sigma|-1}$$

Define

$$\psi(U) = [\mu(X : (\exists \tilde{A} \subseteq U) \varphi(X, \tilde{A})) > 2^{-|\sigma|-1}]$$

The formula $\psi(U)$ is $\Sigma_1^{0,Z}$ and by compactness, $\psi(U)$ is essential. By Z -hyperimmunity of B_i , there exists some finite set $A \subseteq \overline{B_i}$ such that $\psi(A)$ holds. For every set A such that $\psi(A)$ holds, there exists some $X \in \mathcal{S}$ and some $\tilde{A} \subseteq A$ such that $\varphi(X, \tilde{A})$ holds. By definition of $X \in \mathcal{S}$, $\tilde{A} \not\subseteq \overline{B_i}$ and therefore $A \not\subseteq \overline{B_i}$. Contradiction. \square

Proof (Theorem 13). Fix some set Z , some countable collection of Z -hyperimmune sets B_0, B_1, \dots and some Z -computable tree $T \subseteq 2^{<\omega}$. Our forcing conditions are (σ, R) where σ is a stem of the infinite, Z -computable tree $R \subseteq T$. A condition (τ, S) extends (σ, R) if $\sigma \leq \tau$ and $S \subseteq R$. The result is a direct consequence of the following lemma.

Lemma 23. *For every condition $c = (\sigma, R)$, every $\Sigma_1^{0,Z}$ formula $\varphi(G, U)$ and every $i \in \omega$, there exists an extension $d = (\tau, S)$ such that $\varphi(P, U)$ is not essential for every path $P \in [S]$, or $\varphi(\tau, A)$ holds for some $A \subseteq \overline{B_i}$.*

Proof. Define

$$\psi(U) = (\exists s)(\forall \tau \in R \cap 2^s)(\exists \tilde{A} \subseteq_{\text{fin}} U) \varphi(\tau, \tilde{A})$$

The formula $\psi(U)$ is $\Sigma_1^{0,Z}$ so we have two cases:

- Case 1: $\psi(U)$ is not essential with some witness x . By compactness, the following set is an infinite Z -computable subtree of R :

$$S = \{\tau \in R : (\forall A \succ x) \neg \varphi(\tau, A)\}$$

The condition $d = (\sigma, S)$ is an extension such that $\varphi(P, U)$ is not essential for every $P \in [S]$.

- Case 2: $\psi(U)$ is essential. By Z -hyperimmunity of B_i , there exists some finite set $A \subseteq \overline{B_i}$ such that $\psi(A)$ holds. Unfolding the definition of $\psi(A)$, there exists some $\tau \in R$ such that $R^{[\tau]}$ is infinite and $\varphi(\tau, \tilde{A})$ holds for some $\tilde{A} \subseteq A \subseteq \overline{B_i}$. The condition $d = (\tau, R^{[\tau]})$ is an extension such that $\varphi(\tau, \tilde{A})$ holds for some finite set $\tilde{A} \subseteq \overline{B_i}$.

Using Lemma 23, define an infinite descending sequence of conditions $c_0 = (\epsilon, T) \geq c_1 \geq \dots$ such that for each $s \in \omega$

- (i) $|\sigma_s| \geq s$
- (ii) $\varphi(P, U)$ is not essential for every path $P \in [R_{s+1}]$, or $\varphi(\sigma_{s+1}, A)$ holds for some finite set $A \subseteq \overline{B_i}$ if $s = \langle \varphi, i \rangle$

where $c_s = (\sigma_s, R_s)$. \square

C Thin set for pairs and preservation of hyperimmunity

All the proofs in this section are very similar to [21]. We reprove everything in the context of preservation of hyperimmunities for the sake of completeness.

Definition 24 (Strong preservation of n hyperimmunities). A Π_2^1 statement P admits strong preservation of n hyperimmunities if for each set Z , each Z -hyperimmune sets B_0, \dots, B_{n-1} and each (arbitrary) P -instance X , there exists a solution Y to X such that the B 's are $Y \oplus Z$ -hyperimmune.

The following lemma has been proven by the author in its full generality in [17]. We reprove it in the context of preservation of n hyperimmunities.

Lemma 25. For every $k, n \geq 1$ and $m \geq 2$, if TS_m^k admits strong preservation of n hyperimmunities, then TS_m^{k+1} admits preservation of n hyperimmunities.

Proof. Fix any set Z , some Z -hyperimmune sets B_0, \dots, B_{n-1} and any Z -computable coloring $f : [\omega]^{k+1} \rightarrow m$. Consider the uniformly Z -computable sequence of sets $R_{\sigma,i}$ defined for each $\sigma \in [\omega]^k$ and $i < m$ by

$$R_{\sigma,i} = \{s \in \omega : f(\sigma, s) = i\}$$

As COH admits preservation of n hyperimmunities, there exists some \mathbf{R} -cohesive set G such that the B 's are $G \oplus Z$ -hyperimmune. The cohesive set induces a $(G \oplus Z)'$ -computable coloring $\tilde{f} : [\omega]^k \rightarrow m$ defined by:

$$(\forall \sigma \in [\omega]^k) \tilde{f}(\sigma) = \lim_{s \in G} f(\sigma, s)$$

As TS_m^k admits strong preservation of n hyperimmunities, there exists an infinite \tilde{f} -thin set H such that the B 's are $H \oplus G \oplus Z$ -hyperimmune. $H \oplus G \oplus Z$ computes an infinite f -thin set.

Thanks to Lemma 25, it suffices to prove the following theorem to deduce Theorem 21.

Theorem 26. For every $n \geq 1$, TS_{n+1}^1 admits strong preservation of n hyperimmunities.

The remainder of this section is devoted to the proof of Theorem 26. Fix some set Z , some Z -hyperimmune sets B_0, \dots, B_{n-1} and some $(n+1)$ -partition $A_0 \cup \dots \cup A_n = \omega$. We will construct an infinite set G such that $G \cap \bar{A}_i$ is infinite for each $i \leq n$ and the B 's are $(G \cap \bar{A}_i) \oplus Z$ -hyperimmune for some $i \leq n$. Our forcing conditions are Mathias conditions (F, X) such that the B 's are $X \oplus Z$ -hyperimmune.

C.1 Forcing limitlessness

We want to satisfy the following scheme of requirements to ensure that $G \cap \overline{A_i}$ is infinite for each $i \leq n$.

$$\mathcal{Q}_p : (\exists m_0, \dots, m_n > p)[m_0 \in G \cap \overline{A_0} \wedge \dots \wedge m_n \in G \cap \overline{A_n}]$$

We say that an $(n+1)$ -partition $A_0 \cup \dots \cup A_{n-1} = \omega$ is *non-trivial* if there exists no infinite f -thin set H such that the B 's are $H \oplus Z$ -hyperimmune. A condition (F, X) forces \mathcal{Q}_p if there exists $m_0, \dots, m_n > p$ such that $m_i \in F \cap \overline{A_i}$ for each $i \leq n$. Therefore, if G satisfies c and c forces \mathcal{Q}_p , then G satisfies the requirement \mathcal{Q}_p . We now prove that the set of conditions forcing \mathcal{Q}_p is dense for each $p \in \omega$. Therefore, every sufficiently generic filter will induce $n+1$ infinite sets $G \cap \overline{A_0}, \dots, G \cap \overline{A_n}$.

Lemma 27. *For every condition c and every $p \in \omega$, there is a condition d extending c such that d forces \mathcal{Q}_p .*

Proof. Fix some $p \in \omega$. It is sufficient to show that for a condition $c = (F, X)$ and some $i \leq n$, there exists an extension $d_0 = (H, Y)$ and some integer $m_i > p$ that $m_i \in H \cap \overline{A_i}$. By iterating the process for each $i \leq n$, we obtain the desired extension d . Suppose for the sake of contradiction that $X \cap \overline{A_i}$ is finite. Then one can X -compute an infinite set H thin for the A 's with witness j for any $j \neq i$, contradicting non-triviality of f . Therefore, there exists an $m_i \in X \cap \overline{A_i}$, $m_i > p$. The condition $d_0 = (F \cup \{m_i\}, X \setminus [0, m_i])$ is the desired extension. \square

C.2 Forcing non-preservation

Fix an enumeration $\varphi_0(G, U), \varphi_1(G, U), \dots$ of all $\Sigma_1^{0,Z}$ formulas. The second scheme of requirements consists in ensuring that the sets B_0, \dots, B_{n-1} are all $G \cap \overline{A_i}$ -hyperimmune for some $i \leq n$. The requirements are of the following form for each e .

$$\mathcal{R}_e : \bigwedge_{j < n} \mathcal{R}_{e_0}^{A_0, B_j} \vee \dots \vee \bigwedge_{j < n} \mathcal{R}_{e_n}^{A_n, B_j}$$

where

$$\mathcal{R}_e^{A_i, B_j} : \varphi_e(G \cap \overline{A_i}, U) \text{ essential} \Rightarrow (\exists A \subseteq_{\text{fin}} \overline{B_j}) \varphi_e(G \cap \overline{A_i}, A)$$

A condition forces \mathcal{R}_e if every set G satisfying this condition also satisfies the requirement \mathcal{R}_e .

Lemma 28. *For every condition $c = (F, X)$, every $i_0 < i_1 \leq n$, every $j < n$ and every indices e , there exists an extension d such that for some $i \in \{i_0, i_1\}$, d forces $\varphi_e(G \cap \overline{A_i}, U)$ not to be essential or forces $\varphi_e(G \cap \overline{A_i}, A)$ for some finite set $A \subseteq \overline{B_j}$.*

Proof. Let $\psi(U)$ be the formula which holds if for every 2-partition $X_{i_0} \cup X_{i_1} = X$, there is some $i \in \{i_0, i_1\}$ and some set $G_i \subseteq X_i$ such that $\varphi_e((F \cap \overline{A_i}) \cup G_i, \tilde{A})$ holds for some $\tilde{A} \subseteq U$. By compactness, the formula $\psi(U)$ is $\Sigma_1^{0, X \oplus Z}$. We have two cases:

- Case 1: $\psi(U)$ is essential. As B_j is $X \oplus Z$ -hyperimmune, there exists some finite set $A \subseteq \overline{B_j}$ such that $\psi(A)$ holds. In particular, taking $X_i = X \cap \overline{A_i}$ for each $i \in \{i_0, i_1\}$, there exists some $i \in \{i_0, i_1\}$ and some finite set $G_i \subseteq X_i$ such that $\varphi_{e_i}((F \cap \overline{A_i}) \cup G_i, \tilde{A})$ holds for some $\tilde{A} \subseteq A$. The condition $d = (F \cup G_i, X \setminus [0, \max(G_i)])$ is an extension forcing $\varphi_{e_i}(G \cap \overline{A_i}, A)$ for some finite set $A \subseteq \overline{B_j}$.
- Case 2: $\psi(U)$ is not essential, say with witness x . By compactness, the $\Pi_1^{0, X \oplus Z}$ class \mathcal{C} of sets $X_{i_0} \oplus X_{i_1}$ such that $X_{i_0} \cup X_{i_1} = X$ and for every $A > x$, every $i \in \{i_0, i_1\}$ and every set $G_i \subseteq X_i$, $\varphi_{e_i}((F \cap \overline{A_i}) \cup G_i, A)$ does not hold is not empty. By preservation of hyperimmunity of WKL_0 , there exists some $X_{i_0} \oplus X_{i_1} \in \mathcal{C}$ such that the B 's are $X_{i_0} \oplus X_{i_1} \oplus Z$ -hyperimmune. Let $i \in \{i_0, i_1\}$ be such that X_i is infinite. The condition $d = (F, X_i)$ is an extension of c forcing $\varphi_{e_i}(G \cap \overline{A_i}, U)$ not to be essential. \square

Lemma 29. *For every condition c , and every indices e , there exists an extension d forcing \mathcal{R}_e .*

Proof. Fix a condition c , and apply iteratively Lemma 28 to obtain an extension d such that for each $j < n$, d forces $\varphi_{e_i}(G \cap \overline{A_i}, U)$ not to be essential or forces $\varphi_{e_i}(G \cap \overline{A_i}, A)$ for some finite set $A \subseteq \overline{B_j}$ for n different i 's. By the pigeonhole principle, there exists some $i \leq n$ such that d forces $\varphi_{e_i}(G \cap \overline{A_i}, U)$ not to be essential or forces $\varphi_{e_i}(G \cap \overline{A_i}, A)$ for some finite set $A \subseteq \overline{B_j}$ for each $j < n$. Therefore, d forces \mathcal{R}_e . \square

C.3 Construction

Thanks to Lemma 29 and Lemma 27, define an infinite descending sequence of conditions $c_0 = (\emptyset, \omega) \geq c_1 \geq \dots$ such that for each $s \in \omega$,

- (a) c_{s+1} forces \mathcal{R}_e if $s = \langle e \rangle$
- (b) c_{s+1} forces \mathcal{Q}_s

where $c_s = (F_s, X_s)$. Let $G = \bigcup_s F_s$. The sets $G \cap \overline{A_0}, \dots, G \cap \overline{A_n}$ are all infinite and the B 's are $(G \cap \overline{A_i}) \oplus Z$ -hyperimmune for some $i \leq n$. This finishes the proof.