

The strength of the thin set theorems in reverse mathematics

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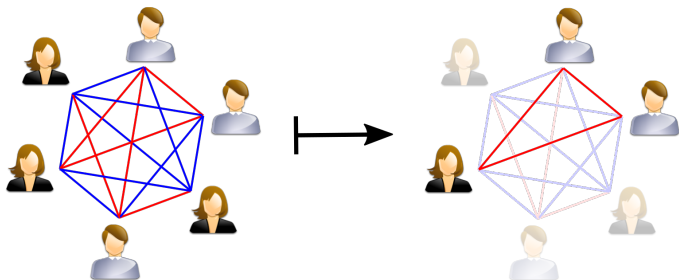
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RAMSEY'S THEOREM

 RT_{k}^n

Every k -coloring of $[\mathbb{N}]^n$ admits an infinite homogeneous set.



REVERSE MATHEMATICS

Q is at least as hard as P if

$$\text{RCA}_0 \vdash Q \rightarrow P$$

in a very weak theory RCA_0
capturing computable mathematics

(Harvey Friedman, 1974)

Turing ideal \mathcal{M}

- ▶ $(\forall X \in \mathcal{M})(\forall Y \leq_T X)[Y \in \mathcal{M}]$
- ▶ $(\forall X, Y \in \mathcal{M})[X \oplus Y \in \mathcal{M}]$

Examples

- ▶ $\{X : X \text{ is computable} \}$
- ▶ $\{X : X \leq_T A \wedge X \leq_T B\}$ for some sets A and B

Let \mathcal{M} be a **Turing ideal** and P, Q be **problems**.

Satisfaction

$$\mathcal{M} \models P$$

if every P -instance in \mathcal{M}
has a solution in \mathcal{M} .

Computable entailment

$$P \models_c Q$$

if every Turing ideal
satisfying P satisfies Q .

Fix two problems P and Q .

How to prove that $P \not\equiv_c Q$?

Build a Turing ideal \mathcal{M} such that

- ▶ $\mathcal{M} \models P$
- ▶ $\mathcal{M} \not\models Q$

A **weakness property** is a collection of sets closed downwards under the Turing reducibility.

Examples

- ▶ $\{X : X \text{ is low}\}$
- ▶ $\{X : A \not\leq_T X\}$ for some set A
- ▶ $\{X : X \text{ is hyperimmune-free}\}$

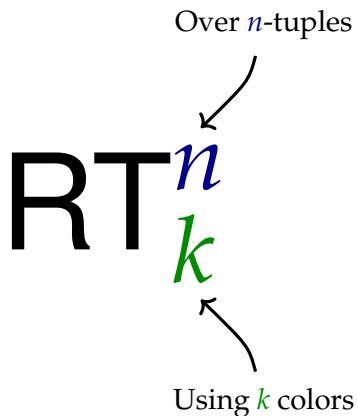
Fix a weakness property \mathcal{W} .

A problem P **preserves** \mathcal{W} if for every $Z \in \mathcal{W}$, every Z -computable P -instance X **has a solution** Y such that $Y \oplus Z \in \mathcal{W}$

Lemma

If P preserves \mathcal{W} but Q does not, then $P \not\equiv_c Q$

RAMSEY'S THEOREM



ACA: $(\forall X)(\exists Y)Y = X'$

ACA \models_c RT_k^n

whenever $n \geq 1$.

(Jockusch, 1972)

$RT_k^n \models_c$ ACA

whenever $n \geq 3$.

(Jockusch, 1972)

$$\text{RT}_k^2 \not\equiv_c \text{ACA}$$

whenever $n, k \geq 1$.

(Seetapun, 1995)

A problem P **avoids cones** if it preserves $\mathcal{W} = \{X : A \not\leq_T X\}$ for every set A .

- ▶ RT_k^2 avoids cones
- ▶ ACA does not avoid cones

WKL: Every infinite **binary tree** has an infinite **path**.

WKL $\not\equiv_c$ **RT**_k²

whenever $k \geq 2$.

(Jockusch, 1972)

RT_k² $\not\equiv_c$ **WKL**

whenever $k \geq 1$.

(Liu, 2012)

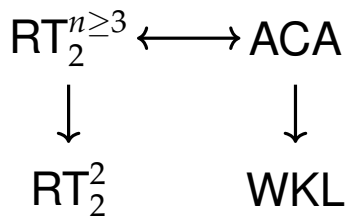
$$\text{RT}_k^2 \not\equiv_c \text{WKL}$$

whenever $k \geq 1$.

(Liu, 2012)

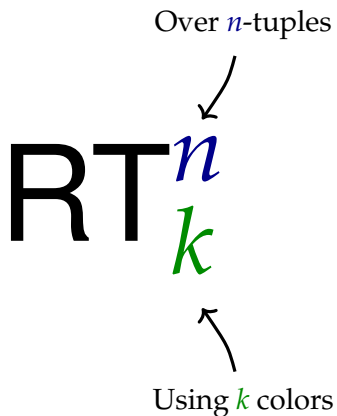
A problem P **avoids PA** if it preserves $\mathcal{W} = \{X : X \text{ is not PA}\}$.

- ▶ RT_k^2 avoids PA
- ▶ WKL does not avoid PA

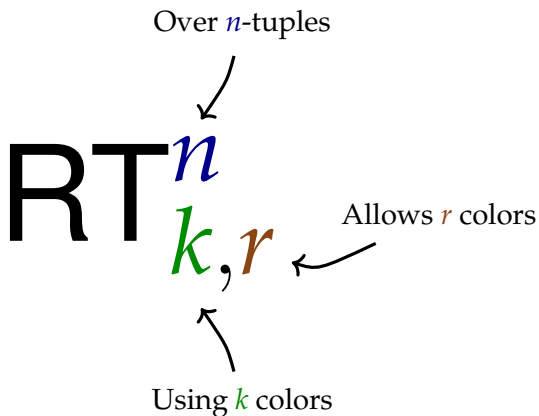


Over computable entailment

RAMSEY'S THEOREM



RAMSEY'S THEOREM



THIN SET THEOREM

 TS_k^n $RT_{k,k-1}^n$

$$\text{TS}_k^n \not\equiv_c \text{ACA}$$

for k is large enough.

(Wang, 2014)

A problem P **avoids cones** if it preserves $\mathcal{W} = \{X : A \not\leq_T X\}$ for every set A .

- ▶ TS_k^n avoids cones when k is sufficiently large
- ▶ ACA does not avoid cones

$$\text{TS}_{k^s}^{ns+1} \models_c \text{TS}_k^{n+1}$$

$$\text{TS}_{2^n}^{n+2} \models_c \text{ACA}$$

(Wang, Dorais, Dzhafarov, Hirst, Mileti, Shafer, 2015)

Tuples	Cone avoidance	Computes \emptyset'
TS_k^2	$k \geq 2$	never
TS_k^3	$k \geq 3$	$k = 2$
TS_k^4	$k \geq 6$	$k \leq 4$

(Jockusch, Wang, Dorais, Dzhafarov,
Hirst, Mileti, Shafer, Cholak, P.)

Does TS_5^4 admit cone avoidance?

$$\text{TS}_{k+1}^2 \not\equiv_c \text{TS}_k^2$$

whenever $k \geq 2$.

(P, 2016)

A problem **P** **preserves k hyperimmunities** if it preserves $\mathcal{W} = \{X : \forall i < k, A_i \text{ is } X\text{-hyperimmune}\}$ for every k -tuple of sets A_0, \dots, A_{k-1} .

- ▶ TS_{k+1}^2 preserves k hyperimmunities
- ▶ TS_k^2 does not preserve k hyperimmunities

$$\text{TS}_\ell^n \not\equiv_c \text{TS}_k^2$$

whenever $\ell \gg k, n$.

(P., 2016)

A problem \mathbf{P} **preserves k hyperimmunities** if it preserves $\mathcal{W} = \{X : \forall i < k, A_i \text{ is } X\text{-hyperimmune}\}$ for every k -tuple of sets A_0, \dots, A_{k-1} .

When $\ell \gg k, n$, TS_ℓ^n preserves k hyperimmunities

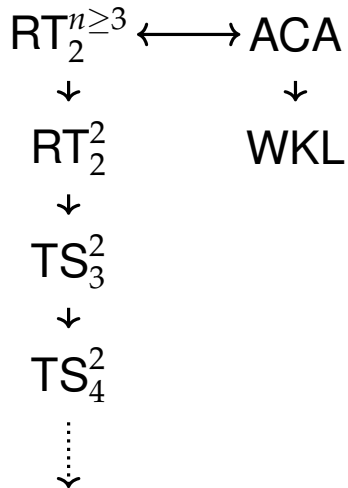
$$\text{TS}_k^n \not\equiv_c \text{WKL}$$

for k is large enough.

(P., 2016)

A problem P **avoids PA** if it preserves $\mathcal{W} = \{X : X \text{ is not PA}\}$.

- ▶ TS_k^n avoids PA when k is sufficiently large
- ▶ WKL does not avoid PA



Over computable entailment

TS^{*n*}

For every coloring $f : [\mathbb{N}]^n \rightarrow \mathbb{N}$,
there is an infinite set H such that

$$f[H]^n \neq \mathbb{N}.$$

FS^{*n*}

For every coloring $f : [\mathbb{N}]^n \rightarrow \mathbb{N}$,
there is an infinite set H such that

$$\forall c \in H, c \notin f[H \setminus \{c\}]^n.$$

$$\text{TS}_k^n \models_c \text{TS}^n$$

$$\text{FS}^n \models_c \text{TS}^n$$

Is FS^n **strictly stronger** than TS^n ?

$$\text{TS}_3^2 \not\equiv_c \text{FS}^2$$

(P., 2017)

A formula $\varphi(\vec{U})$ is **essential** if for every $x \in \omega$, there are some sets $\vec{V} > x$ such that $\varphi(\vec{V})$ holds.

A function $f : \omega \rightarrow \omega$ is **freely X -hyperimmune** if for every $n \in \omega$, every essential $\Sigma_1^{0,X}$ formula $\varphi(\vec{U})$ and every function $g : \mathcal{P}_+(\omega) \rightarrow \omega$, $\varphi(\vec{V})$ holds for some sets \vec{V} such that

$$(\forall y \in \bigcup_i V_i) f(y) = g(\{i : y \in V_i\})$$

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